KERNEL VERBS API UPDATE
Leon Romanovsky
Mellanox Technologies
[ April 4th, 2016 ]
AGENDA

- Memory registration API
- CQ polling API
- Draining QP
- Generic RDMA WRITE/READ API
MEMORY REGISTRATION API
MULTIPLE MEMORY REGISTRATION METHODS

- **Physical memory regions (MR)**
  - Synchronous interface
  - Every registration causes a new MR

- **Fast memory registration (FMR)**
  - Fast synchronous interface
  - Weak deregistration semantics
  - Not widely adopted

- **Memory windows (MW)**
  - Fast, asynchronous interface
  - Binds continuous apertures to existing memory regions
    - Not relevant to kernel ULPs
    - Removed from kernel API

- **Fast memory registration mode (FRWR)**
  - Asynchronous interface
  - Maps blocks of physical memory
  - Widely adopted
MULTIPLE MEMORY REGISTRATION METHODS

- **Physical memory regions (MR)**
  - Synchronous interface
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- **Fast memory registration mode (FRWR)**
  - Asynchronous interface
  - Maps blocks of physical memory
  - Widely adopted
- All buffers must be block aligned (page shift)
- The first buffer can have first byte offset (FBO)
- The last buffer can end before the end of block (EOB)
FRWR DRAWBACKS

- Allocation of free memory region and a fast_reg_page_list
- Translation from S/G list to page vector for each ULP
- No ability to support arbitrary S/G
  - Each ULP bridged this semantic gap
## OLD FRWR API

<table>
<thead>
<tr>
<th>OP</th>
<th>API</th>
</tr>
</thead>
<tbody>
<tr>
<td>Allocate</td>
<td><code>frpl = ib_alloc_fast_reg_page_list()</code></td>
</tr>
<tr>
<td></td>
<td><code>mr = ib_alloc_fast_reg_mr()</code></td>
</tr>
<tr>
<td>Free</td>
<td><code>ib_free_fast_reg_page_list(frpl)</code></td>
</tr>
<tr>
<td></td>
<td><code>ib_dereg_mr(mr)</code></td>
</tr>
<tr>
<td>Register interface</td>
<td><code>struct {</code></td>
</tr>
<tr>
<td>(post_send)</td>
<td><code>u64</code></td>
</tr>
<tr>
<td></td>
<td><code>struct ib_fast_reg_page_list *page_list;</code></td>
</tr>
<tr>
<td></td>
<td><code>unsigned int</code></td>
</tr>
<tr>
<td></td>
<td><code>unsigned int</code></td>
</tr>
<tr>
<td></td>
<td><code>u32</code></td>
</tr>
<tr>
<td></td>
<td><code>int</code></td>
</tr>
<tr>
<td></td>
<td><code>u32</code></td>
</tr>
<tr>
<td></td>
<td>} fast_reg;</td>
</tr>
<tr>
<td></td>
<td><code>iova_start;</code></td>
</tr>
<tr>
<td></td>
<td><code>*page_list;</code></td>
</tr>
<tr>
<td></td>
<td><code>page_shift;</code></td>
</tr>
<tr>
<td></td>
<td><code>page_list_len;</code></td>
</tr>
<tr>
<td></td>
<td><code>length;</code></td>
</tr>
<tr>
<td></td>
<td><code>access_flags;</code></td>
</tr>
<tr>
<td></td>
<td><code>rkey;</code></td>
</tr>
<tr>
<td>Opcode</td>
<td>IB_WR_FAST_REG_MR</td>
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</table>
## NEW FRWR API

<table>
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<tr>
<td><strong>Allocate</strong></td>
<td>mr = ib_alloc_mr(pd, max_num_sg, mr_type)</td>
</tr>
<tr>
<td><strong>Free</strong></td>
<td>ib_dereg_mr(mr)</td>
</tr>
<tr>
<td><strong>S/G list mapping</strong></td>
<td>nents = ib_map_map_mr_sg(mr, sg, sg_nents, page_size);</td>
</tr>
</tbody>
</table>
| **Register interface (post_send)** | struct {  
            struct ib_send_wr                      wr;  
            struct ib_mr                          *mr;  
            int                                    access;  
            u32                                    key;  
        } ib_reg_mr;                                |
| **Opcode**          | IB_WR_REG_MR                                                        |
ULP REGISTRATION FLOW

Allocate MR

Post send – reg MR

Post send – transfer key
CQ POLLING API
CQ POLLING CONSIDERATIONS

- Different completion contexts
  - Kernel threads
  - Work queues
  - Software IRQs
  - Hardware IRQs
- Fairness between multiple CQs
- CQ re-arm policy
- Handling missed events
CQ POLLING USAGE UNCERTAINTY

- Work requests (WR) context return path
- Work request IDs (wr_id) (un)reliable
- Multiple sources of post_send completions
- Polling in batches or one-by-one
- Multiple CPUs (affinity)
- CPU/NUMA locality
## CQ POLLING API

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| Decide on polling type  | ```
enum ib_poll_context {
    IB_POLL_DIRECT, /* caller context, no hw completions */
    IB_POLL_SOFTIRQ, /* poll from softirq context */
    IB_POLL_WORKQUEUE, /* poll from workqueue */
};
``` |
| Allocate               | ```
struct ib_cq *ib_alloc_cq(dev, private, nr_cqe, comp_vector, poll_ctx);
``` |
| Free                   | ```
void ib_free_cq(cq)
``` |
| Completion return      | ```
struct ib_cqe {
    void (*done)(struct ib_cq *cq, struct ib_wc *wc);
};
``` |
| (filled in WR)          |                                                                     |
CQ POLLING SUMMARY

- Unify completion queue logic from different ULP clients
- Simplify completion queue polling and interrupt handling
  - No need to poll, handle events and maintain logic
- Resolve the error completions unreliability
- Support different polling schemes
- Performance optimized
DRAINING QP
PROBLEM

- Unknown when WRs are completed after ceasing posts
1. Wait until all previously posted WRs complete and then destroy QP - may be indefinite

2. Destroy QP without waiting for completions - need to maintain a shadow state of all in-flight WQEs in order to free related state

3. Modify QP to error, and then poll related CQ until it is empty - works only if CQ is associated with only the same QP
ROBUST GENERIC DRAINING

1. Cease posting new WRs to QP
2. Change QP state to ERR
3. Post marker WR (nop)
4. Wait until marker WR completes
### DRAIN QP API

<table>
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<tr>
<td>Drain SQ</td>
<td><code>void ib_drain_sq(qp)</code></td>
</tr>
<tr>
<td>Drain RQ</td>
<td><code>void ib_drain_rq(qp)</code></td>
</tr>
<tr>
<td>Drain QP</td>
<td><code>void ib_drain_qp(qp)</code></td>
</tr>
</tbody>
</table>

- **IB/core code to drain SQ/RQ/both in single function call**
- **Protect from use-after-free error flow**
- **Synchronized operation in one place**
GENERIC RDMA READ/WRITE (WIP)
MOTIVATION

- RDMA READ/WRITE is common operation for storage protocols
- Every ULP has own implementation
- Lack of support for generic S/G lists
- HCA aware implementations
- Reuse pre-allocated MRs (MR pool)
- Support large number of S/G entries
REFERENCES

- **IB: new common API for draining queues** by Steve Wise
- **Generic RDMA READ/WRITE API** by Christoph Hellwig
- **Completion queue abstraction** by Christoph Hellwig and Sagi Grimberg
- **New fast registration API** by Sagi Grimberg
12th ANNUAL WORKSHOP 2016

THANK YOU

Leon Romanovsky
Mellanox Technology